# CSCB63 - Design and Analysis of Data Structures 

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## how long do things take

Remember this algorithm?

| 1 | $i=1$ |
| :--- | :--- |
| 2 | while $i<\operatorname{len}(A):$ |
| 3 | $v=A[i]$ |
| 4 | $j=i$ |
| 5 | while $j>0$ and $A[j-1]>v:$ |
| 6 | $A[j]=A[j-1]$ |
| 7 | $j=j-1$ |
| 8 | $A[j]=v$ |
| 9 | $i=i+1$ |

What do we count? Does it matter?

## how long do things take

Let's try counting this way:

- get/set variables: 1 step
- function call: $1+$ steps to evaluate each argument + steps to execute function
- return statement: $1+$ steps to evaluate return value
- if/while condition: $1+$ steps to evaluate the boolean expression
- assignment statement: $1+$ steps to evaluate each side
- arithmetic/comparison/boolean operators: $1+$ steps to evaluate each operand
- array access: $1+$ steps to evaluate array index
- constants: free!


## how long do things take

| 0 def | InsertionSort (A) : | STEPS |
| :---: | :---: | :---: |
| 1 | i $=1$ | 2 |
| 2 | while i < len(A): | 5 |
| 3 | $\mathrm{v}=\mathrm{A}$ [i] | 5 |
| 4 | j $=$ i | 3 |
| 5 | while $\mathrm{j}>0$ and $\mathrm{A}[\mathrm{j}-1]$ > v: | 10 or 3 |
| 6 | $\mathrm{A}[\mathrm{j}]=\mathrm{A}[\mathrm{j}-1]$ | 8 |
| 7 | $\mathrm{j}=\mathrm{j}-1$ | 4 |
| 8 | $\mathrm{A}[\mathrm{j}]=\mathrm{v}$ | 5 |
| 9 | $\mathrm{i}=\mathrm{i}+1$ | 4 |

What assumptions did we make? Are they realistic?

So, what's the total number of steps?

## how long do things take

In the worst case:

- line 1: once: 2 steps

For $n \geq 1$ :

- line 2: $n-1$ times (true) +1 time (false): $5 n$ steps
- lines $3,4,8,9: n-1$ times : $(5+3+5+4)(n-1)=17 n-17$ steps
- line 5: for each $i$ : $i$ times (true) +1 time (false) : $10 i+3$ steps
- lines 6, 7: for each $i: i$ times : $(8+4) i=12 i$ steps

$$
\begin{aligned}
& 2+5 n+17 n-17+\sum_{i=1}^{n-1}(10 i+3+ \\
= & 22 n-15+\sum_{i=1}^{n-1}(22 i+3) \\
= & 22 n-15+22 \frac{(n-1) n}{2}+3(n-1) \\
= & 11 n^{2}+14 n-18
\end{aligned}
$$

## how long do things take

In the best case:

- line 1: once : 2 steps

For $n \geq 1$ :

- line 2: $n-1$ times (true) +1 time (false) : $5 n$ steps
- lines $3,4,8,9: n-1$ times :

$$
(5+3+5+4)(n-1)=17 n-17 \text { steps }
$$

- line 5: for each $i: 1$ time (false) : 10 steps
- lines 6, 7: for each $i$ : 0 times: 0 steps

$$
\begin{aligned}
2+5 n+17 n-17+\sum_{i=1}^{n-1} 10 & =22 n-15+(n-1) 10 \\
& =32 n-25
\end{aligned}
$$

## how long do things take

What if we write the same algorithm differently?
0 def InsertionSort (A):
$1 \quad \mathrm{n}=\operatorname{len}(\mathrm{A})$
2 for (i = 1; i < n; i++):
3 for (j = i; j > 0 and $A[j]<A[j-1] ; j--):$
4 swap A[j], A[j-1]

- line 1: once, 4 steps

For $n \geq 1$ :

- line 2: 1 step (once) +2 steps (once) +3 steps ( $n$ times) + 2 steps ( $n-1$ times)
- line 3: for each $i$ : 1 step (once) +3 steps (once) +11 steps ( $i$ times) +2 steps (once) +2 steps ( $i$ times)
- line 4: for each $i$ : 9 steps ( $i$ times)


## how long do things take

- line 1: once, 4 steps

For $n \geq 1$ :

- line 2: 1 step (once) +2 steps (once) +3 steps ( $n$ times) + 2 steps ( $n-1$ times)
- line 3: for each $i$ : 1 step (once) +3 steps (once) +11 steps ( $i$ times) +2 steps (once) +2 steps ( $i$ times)
- line 4: for each $i: 9$ steps ( $i$ times)

$$
\begin{aligned}
& 4+1+2+3 n+2(n-1)+\sum_{i=1}^{n-1}(1+3+11 i+2+2 i+9 i) \\
= & 5 n+5+\sum_{i=1}^{n-1}(22 i+6) \\
= & 5 n+5+22 \frac{(n-1) n}{2}+6(n-1) \\
= & 11 n^{2}-1
\end{aligned}
$$

Is this the same running time? In what sense?

## how long do things take

Q. What if I now run this algorithm on a machine that is slower to perform variable look up and write?
Q. Should the complexity change?
Q. How important are those constants as the input size $n$ gets large?
Q. How are our two results $11 n^{2}+14 n-18$ and $11 n^{2}-1$ similar?
Q. They are both quadratic polynomials.

## how long do things take

We say...

- that a quadratic polynomial is of order $n^{2}$,
- that a cubic polynomial is of order $n^{3}$,
- that $4 n \lg (n)+2 n+10$ is of order $n \lg (n)$.

Why can we say this? With a little mathemagic:

$$
11 n^{2}+14 n-18 \leq 11 n^{2}+14 n \leq 11 n^{2}+14 n^{2}=25 n^{2}
$$

Another example:

$$
11 n^{2}-21 n+19 \leq 11 n^{2}+19 \leq 11 n^{2}+n \leq 11 n^{2}+n^{2}=12 n^{2}
$$

for all natural $n \geq 19$

## how long do things take - formally

For all natural $n \geq 19$ :

$$
11 n^{2}-21 n+19 \leq 12 n^{2}
$$

There exists an $n_{0} \in \mathbb{N}$ such that, for all natural $n \geq n_{0}$,

$$
11 n^{2}-21 n+19 \leq 12 n^{2}
$$

We can take this even further and say, there exists real $c>0$ and natural $n_{0}$ such that, for all natural $n \geq n_{0}$,

$$
11 n^{2}-21 n+19 \leq c \cdot n^{2}
$$

which is exactly the definition of "Big-Oh"!

## Big-Oh - Asymptotic Upper Bound

We denote:

- $\mathbb{N}$ : the set of natural numbers
- $\mathbb{R}^{+}$: the set of positive real numbers
- $\mathcal{F}$ : the set of functions $f: \mathbb{N} \rightarrow \mathbb{R}^{+}$

Let $g \in \mathcal{F}$. Define $\mathcal{O}(g)$ to be the set of functions $f \in \mathcal{F}$ such that

$$
\exists c \in \mathbb{R}^{+}, \exists n_{0} \in \mathbb{N}, \forall n \in \mathbb{N}, n \geq n_{0} \Rightarrow f(n) \leq c \cdot g(n)
$$

