CSCB63 – Design and Analysis of Data Structures

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¹based on notes by Anna Bretscher and Albert Lai

Remember this algorithm?

```
1
      i = 1
2
      while i < len(A):
3
        v = A[i]
4
        j = i
5
        while j > 0 and A[j-1] > v:
6
          A[j] = A[j-1]
7
          j = j - 1
        A[j] = v
8
        i = i + 1
9
```

What do we count? Does it matter?

Let's try counting this way:

- get/set variables: 1 step
- function call: 1 + steps to evaluate each argument + steps to execute function
- return statement: 1 + steps to evaluate return value
- if/while condition: 1 + steps to evaluate the boolean expression
- assignment statement: 1 + steps to evaluate each side
- arithmetic/comparison/boolean operators: 1 + steps to evaluate each operand
- array access: 1 + steps to evaluate array index
- constants: free!

| 0 | def | InsertionSort (A): | STEPS |
|---|-----|----------------------------------|---------|
| 1 | | i = 1 | 2 |
| 2 | | while i < len(A): | 5 |
| 3 | | v = A[i] | 5 |
| 4 | | j = i | 3 |
| 5 | | while $j > 0$ and $A[j-1] > v$: | 10 or 3 |
| 6 | | A[j] = A[j-1] | 8 |
| 7 | | j = j - 1 | 4 |
| 8 | | A[j] = v | 5 |
| 9 | | i = i + 1 | 4 |

What assumptions did we make? Are they realistic?

So, what's the total number of steps?

In the worst case:

• line 1: once : 2 steps

For $n \geq 1$:

- line 2: n-1 times (true) + 1 time (false) : 5n steps
- lines 3, 4, 8, 9: n-1 times : (5+3+5+4)(n-1) = 17n-17 steps
- line 5: for each *i*: *i* times (true) + 1 time (false) : 10i + 3 steps

• lines 6, 7: for each *i*: *i* times :
$$(8 + 4)i = 12i$$
 steps
 $2 + 5n + 17n - 17 + \sum_{i=1}^{n-1} (10i + 3 + 12i)$
 $= 22n - 15 + \sum_{i=1}^{n-1} (22i + 3)$
 $= 22n - 15 + 22 \frac{(n-1)n}{2} + 3(n-1)$
 $= 11n^2 + 14n - 18$

In the best case:

• line 1: once : 2 steps

For $n \geq 1$:

• line 2: n-1 times (true) + 1 time (false) : 5n steps

• lines 3, 4, 8, 9:
$$n-1$$
 times :
 $(5+3+5+4)(n-1) = 17n - 17$ steps

- line 5: for each i: 1 time (false) : 10 steps
- lines 6, 7: for each *i*: 0 times : 0 steps

$$2 + 5n + 17n - 17 + \sum_{i=1}^{n-1} 10 = 22n - 15 + (n-1)10$$
$$= 32n - 25$$

What if we write the same algorithm differently?

| 0 | def | InsertionSort (A): |
|---|-----|---|
| 1 | | n = len(A) |
| 2 | | for $(i = 1; i < n; i++):$ |
| 3 | | for $(j = i; j > 0 \text{ and } A[j] < A[j-1]; j):$ |
| 4 | | swap A[j], A[j-1] |

• line 1: once, 4 steps

For $n \geq 1$:

- line 2: 1 step (once) + 2 steps (once) + 3 steps (n times) + 2 steps (n 1 times)
- line 3: for each i: 1 step (once) + 3 steps (once) + 11 steps (i times) + 2 steps (once) + 2 steps (i times)
- line 4: for each *i*: 9 steps (*i* times)

• line 1: once, 4 steps

For $n \geq 1$:

- line 2: 1 step (once) + 2 steps (once) + 3 steps (n times) + 2 steps (n 1 times)
- line 3: for each i: 1 step (once) + 3 steps (once) + 11 steps (i times) + 2 steps (once) + 2 steps (i times)
- line 4: for each *i*: 9 steps (*i* times)

$$4 + 1 + 2 + 3n + 2(n - 1) + \sum_{i=1}^{n-1} (1 + 3 + 11i + 2 + 2i + 9i)$$

$$=5n + 5 + \sum_{i=1}^{n-1} (22i + 6)$$
$$=5n + 5 + 22 \frac{(n-1)n}{2} + 6(n-1)$$
$$=11n^2 - 1$$

Is this the same running time? In what sense?

Q. What if I now run this algorithm on a machine that is slower to perform variable look up and write?

Q. Should the complexity change?

Q. How important are those constants as the input size *n* gets large?

Q. How are our two results $11n^2 + 14n - 18$ and $11n^2 - 1$ similar?

Q. They are both quadratic polynomials.

We say...

- that a quadratic polynomial is of order n^2 ,
- that a cubic polynomial is of order n^3 ,
- that $4n \lg(n) + 2n + 10$ is of order $n \lg(n)$.

Why can we say this? With a little mathemagic:

$$11n^2 + 14n - 18 \le 11n^2 + 14n \le 11n^2 + 14n^2 = 25n^2$$

Another example:

$$11n^2 - 21n + 19 \le 11n^2 + 19 \le 11n^2 + n \le 11n^2 + n^2 = 12n^2$$

for all natural $n > 19$

how long do things take — formally For all natural $n \ge 19$:

$$11n^2 - 21n + 19 \le 12n^2$$

There exists an $n_0 \in \mathbb{N}$ such that, for all natural $n \ge n_0$,

$$11n^2 - 21n + 19 \le 12n^2$$

We can take this even further and say, there exists real c > 0 and natural n_0 such that, for all natural $n \ge n_0$,

$$11n^2 - 21n + 19 \le c \cdot n^2$$

which is exactly the definition of "Big-Oh"!

Big-Oh — Asymptotic Upper Bound

We denote:

- \mathbb{N} : the set of natural numbers
- \mathbb{R}^+ : the set of positive real numbers
- \mathcal{F} : the set of functions $f : \mathbb{N} \to \mathbb{R}^+$

Let $g \in \mathcal{F}$. Define $\mathcal{O}(g)$ to be the set of functions $f \in \mathcal{F}$ such that $\exists c \in \mathbb{R}^+, \exists n_0 \in \mathbb{N}, \forall n \in \mathbb{N}, n \ge n_0 \Rightarrow f(n) \le c \cdot g(n)$