

# CSCB20 – Week 3

## *Introduction to Database and Web Application Programming*

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# This Week

Intro to SQL and MySQL

Mapping Relational Algebra to SQL queries

We will focus on queries to start – assume tables and database exist.

Time permitting:

creating tables, more involved queries...

# Projection

Symbol is  $\Pi$

Selection of attributes.

$\Pi_{ID, salary}(instructor)$

SQL Notation:

SELECT col\_1, ..., col\_N FROM instructor

Or

SELECT \* FROM instructor (means select all columns)

| ID    | name       | dept_name  | salary |
|-------|------------|------------|--------|
| 10101 | Srinivasan | Comp. Sci. | 65000  |
| 12121 | Wu         | Finance    | 90000  |
| 15151 | Mozart     | Music      | 40000  |
| 22222 | Einstein   | Physics    | 95000  |
| 32343 | El Said    | History    | 60000  |
| 33456 | Gold       | Physics    | 87000  |
| 45565 | Katz       | Comp. Sci. | 75000  |
| 58583 | Califieri  | History    | 62000  |
| 76543 | Singh      | Finance    | 80000  |
| 76766 | Crick      | Biology    | 72000  |
| 83821 | Brandt     | Comp. Sci. | 92000  |
| 98345 | Kim        | Elec. Eng. | 80000  |

SELECT ID, salary FROM instructor

# Selection

Notation is  $\sigma_p(x)$ .

$\sigma_{\text{salary} \geq 85000}(\text{instructor})$

SQL Notation:

SELECT \* FROM instructor WHERE salary >= 85000

SELECT col\_1, ..., col\_N FROM instructor WHERE salary >= 85000

| ID    | name       | dept_name  | salary |
|-------|------------|------------|--------|
| 10101 | Srinivasan | Comp. Sci. | 65000  |
| 12121 | Wu         | Finance    | 90000  |
| 15151 | Mozart     | Music      | 40000  |
| 22222 | Einstein   | Physics    | 95000  |
| 32343 | El Said    | History    | 60000  |
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| 76543 | Singh      | Finance    | 80000  |
| 76766 | Crick      | Biology    | 72000  |
| 83821 | Brandt     | Comp. Sci. | 92000  |
| 98345 | Kim        | Elec. Eng. | 80000  |

# Natural Join

Recall we combine two relations into a single relation.

The tuples are joined if the attributes common to both relations are equal.

instructor

| <i>ID</i> | <i>name</i> | <i>dept_name</i> | <i>salary</i> |
|-----------|-------------|------------------|---------------|
| 10101     | Srinivasan  | Comp. Sci.       | 65000         |
| 12121     | Wu          | Finance          | 90000         |
| 15151     | Mozart      | Music            | 40000         |
| 22222     | Einstein    | Physics          | 95000         |
| 32343     | El Said     | History          | 60000         |
| 33456     | Gold        | Physics          | 87000         |
| 45565     | Katz        | Comp. Sci.       | 75000         |
| 58583     | Califieri   | History          | 62000         |
| 76543     | Singh       | Finance          | 80000         |
| 76766     | Crick       | Biology          | 72000         |
| 83821     | Brandt      | Comp. Sci.       | 92000         |
| 98345     | Kim         | Elec. Eng.       | 80000         |



department

| <i>dept_name</i> | <i>building</i> | <i>budget</i> |
|------------------|-----------------|---------------|
| Biology          | Watson          | 90000         |
| Comp. Sci.       | Taylor          | 100000        |
| Elec. Eng.       | Taylor          | 85000         |
| Finance          | Painter         | 120000        |
| History          | Painter         | 50000         |
| Music            | Packard         | 80000         |
| Physics          | Watson          | 70000         |

# Natural Join

instructor ⋈ department

The tuples are joined if the attributes common to both relations are equal.

| <i>ID</i> | <i>name</i> | <i>salary</i> | <i>dept_name</i> | <i>building</i> | <i>budget</i> |
|-----------|-------------|---------------|------------------|-----------------|---------------|
| 10101     | Srinivasan  | 65000         | Comp. Sci.       | Taylor          | 100000        |
| 12121     | Wu          | 90000         | Finance          | Painter         | 120000        |
| 15151     | Mozart      | 40000         | Music            | Packard         | 80000         |
| 22222     | Einstein    | 95000         | Physics          | Watson          | 70000         |
| 32343     | El Said     | 60000         | History          | Painter         | 50000         |
| 33456     | Gold        | 87000         | Physics          | Watson          | 70000         |
| 45565     | Katz        | 75000         | Comp. Sci.       | Taylor          | 100000        |
| 58583     | Califieri   | 62000         | History          | Painter         | 50000         |
| 76543     | Singh       | 80000         | Finance          | Painter         | 120000        |
| 76766     | Crick       | 72000         | Biology          | Watson          | 90000         |
| 83821     | Brandt      | 92000         | Comp. Sci.       | Taylor          | 100000        |
| 98345     | Kim         | 80000         | Elec. Eng.       | Taylor          | 85000         |

SQL Notation:

SELECT \* FROM instructor NATURAL JOIN department

# Cartesian Product Example

Relations  $r, s$ :

| A        | B |
|----------|---|
| $\alpha$ | 1 |
| $\beta$  | 2 |

$r$

| C        | D  | E |
|----------|----|---|
| $\alpha$ | 10 | a |
| $\beta$  | 10 | a |
| $\beta$  | 20 | b |
| $\gamma$ | 10 | b |

$s$

$r \times s$ :

SQL Notation:

SELECT \* FROM  $r$  INNER JOIN  $s$

or

SELECT \* FROM  $r, s$

| A        | B | C        | D  | E |
|----------|---|----------|----|---|
| $\alpha$ | 1 | $\alpha$ | 10 | a |
| $\alpha$ | 1 | $\beta$  | 10 | a |
| $\alpha$ | 1 | $\beta$  | 20 | b |
| $\alpha$ | 1 | $\gamma$ | 10 | b |
| $\beta$  | 2 | $\alpha$ | 10 | a |
| $\beta$  | 2 | $\beta$  | 10 | a |
| $\beta$  | 2 | $\beta$  | 20 | b |
| $\beta$  | 2 | $\gamma$ | 10 | b |

Note: can have as many relations as needed...but what may be a concern?

# Cartesian Product Example

Relations  $r$ ,  $s$ :

$r \times s$ :

| A        | B |
|----------|---|
| $\alpha$ | 1 |
| $\beta$  | 2 |

$r$

| C        | D  | E |
|----------|----|---|
| $\alpha$ | 10 | a |
| $\beta$  | 10 | a |
| $\beta$  | 20 | b |
| $\gamma$ | 10 | b |

$s$

SQL Notation:

SELECT \* FROM  $r$  INNER JOIN  $s$

What if we don't want ALL rows?

For example, we want rows where  $A$ 's value and  $C$ 's value are equal?

| A        | B | C        | D  | E |
|----------|---|----------|----|---|
| $\alpha$ | 1 | $\alpha$ | 10 | a |
| $\alpha$ | 1 | $\beta$  | 10 | a |
| $\alpha$ | 1 | $\beta$  | 20 | b |
| $\alpha$ | 1 | $\gamma$ | 10 | b |
| $\beta$  | 2 | $\alpha$ | 10 | a |
| $\beta$  | 2 | $\beta$  | 10 | a |
| $\beta$  | 2 | $\beta$  | 20 | b |
| $\beta$  | 2 | $\gamma$ | 10 | b |

SELECT \* FROM  $r$  INNER JOIN  $s$  ON  $A = C$



# Inner Join

SQL Notation:

```
SELECT Column1, Column2, ..., ColumnK FROM  
    TableA INNER JOIN TableB  
    ON join_constraints  
    WHERE constraints  
    ORDER BY ColumnX
```

There are many other options, we will see these later...

# Self Join

Suppose we want to join a table to itself.

We want to find those departments that are in the same building.

department A

| <i>dept_name</i> | <i>building</i> | <i>budget</i> |
|------------------|-----------------|---------------|
| Biology          | Watson          | 90000         |
| Comp. Sci.       | Taylor          | 100000        |
| Elec. Eng.       | Taylor          | 85000         |
| Finance          | Painter         | 120000        |
| History          | Painter         | 50000         |
| Music            | Packard         | 80000         |
| Physics          | Watson          | 70000         |

department B

| <i>dept_name</i> | <i>building</i> | <i>budget</i> |
|------------------|-----------------|---------------|
| Biology          | Watson          | 90000         |
| Comp. Sci.       | Taylor          | 100000        |
| Elec. Eng.       | Taylor          | 85000         |
| Finance          | Painter         | 120000        |
| History          | Painter         | 50000         |
| Music            | Packard         | 80000         |
| Physics          | Watson          | 70000         |

SQL Notation:

```
SELECT A.dept_name, B.dept_name FROM
    department A INNER JOIN department B
    ON A.building = B.building
```

# Union

Relations  $r, s$ :

For  $r \cup s$  to be valid.

1.  $r, s$  must have the *same arity* (same number of attributes)
2. The attribute domains must be *compatible*

MySQL Notation:

```
(SELECT * FROM r)
UNION
(SELECT * FROM s)
```

Use UNION ALL to keep duplicates.

| A        | B |
|----------|---|
| $\alpha$ | 1 |
| $\alpha$ | 2 |
| $\beta$  | 1 |

$r$

| A        | B |
|----------|---|
| $\alpha$ | 2 |
| $\beta$  | 3 |

$s$

| A        | B |
|----------|---|
| $\alpha$ | 1 |
| $\alpha$ | 2 |
| $\beta$  | 1 |
| $\beta$  | 3 |

$r \cup s$ :

# Intersection

Relation  $r, s$ :

$$r \cap s = r - (r - s)$$

| $A$      | $B$ |
|----------|-----|
| $\alpha$ | 1   |
| $\alpha$ | 2   |
| $\beta$  | 1   |

$r$

| $A$      | $B$ |
|----------|-----|
| $\alpha$ | 2   |
| $\beta$  | 3   |

$s$

SQL Notation:

(SELECT \* FROM  $r$ )  
INTERSECT  
(SELECT \* FROM  $s$ )

| $A$      | $B$ |
|----------|-----|
| $\alpha$ | 2   |

# Intersection

SQL Notation:

(SELECT \* FROM r)

INTERSECT Does NOT Work in MySQL

(SELECT \* FROM s)

| A        | B |
|----------|---|
| $\alpha$ | 1 |
| $\alpha$ | 2 |
| $\beta$  | 1 |

*r*

| A        | B |
|----------|---|
| $\alpha$ | 2 |
| $\beta$  | 3 |

*s*

MySQL Options:

| A        | B |
|----------|---|
| $\alpha$ | 2 |

LEFT AS EXERCISE

# Difference

What would you expect them to be?

Relations  $r$ ,  $s$ :

$r - s$

SQL Notation:

(SELECT \* FROM  $r$ )  
EXCEPT  
(SELECT \* FROM  $s$ )

| $A$      | $B$ |
|----------|-----|
| $\alpha$ | 1   |
| $\alpha$ | 2   |
| $\beta$  | 1   |

$r$

| $A$      | $B$ |
|----------|-----|
| $\alpha$ | 2   |
| $\beta$  | 3   |

$s$

| $A$      | $B$ |
|----------|-----|
| $\alpha$ | 1   |
| $\beta$  | 1   |

# Difference

SQL Notation:

(SELECT \* FROM r)

EXCEPT Does NOT Work in MySQL

(SELECT \* FROM s)

MySQL Options:

LEFT AS EXERCISE

| A        | B |
|----------|---|
| $\alpha$ | 1 |
| $\alpha$ | 2 |
| $\beta$  | 1 |

*r*

| A        | B |
|----------|---|
| $\alpha$ | 2 |
| $\beta$  | 3 |

*s*

| A        | B |
|----------|---|
| $\alpha$ | 1 |
| $\beta$  | 1 |

# SQL Types

`char( $n$ )`: A fixed-length character string.

`varchar( $n$ )`: A variable-length character string with max length  $n$ .

`int`: An integer.

`numeric( $p, d$ )`: A fixed-point number with  $p$  digits of which  $d$  of the digits are to the right of the decimal point.

`real, double precision`: Floating point and double precision floating point.

`float( $n$ )`: A floating point with at least  $n$  digits of precision.



# Null Value

Every type can have the special value null.

A value of null indicates the value is unknown or that it may not exist at all.

Sometimes we do not want a null value at all – we can add such a constraint.

# Creating a Table

SQL Notation:

```
CREATE TABLE table_name  
    (col_name1 type1,  
    col_name2 type2,  
    ... ,  
    col_namen typen,  
    <integrity-constraint1>,  
    ... ,  
    <integrity-constraintk>);
```

# Integrity Constraints

Primary key( list of attributes) :

These attributes form the primary keys for the relation. Primary keys must be *non-null* and *unique*.

Foreign key( list of attributes) references s :

The values of these attributes for any tuple in the relation must correspond to values of the *primary key attributes* of some tuple in relation s.

not null:

Specifies that this attribute may not have the *null value*. We list this constraint when defining the type of the attribute.

# Examples

CREATE TABLE department

```
(dept_name    VARCHAR(20),  
 building     VARCHAR(15),  
 budget       NUMERIC(12,2),  
 PRIMARY KEY (dept_name));
```

| <i>dept_name</i> | <i>building</i> | <i>budget</i> |
|------------------|-----------------|---------------|
| Biology          | Watson          | 90000         |
| Comp. Sci.       | Taylor          | 100000        |
| Elec. Eng.       | Taylor          | 85000         |
| Finance          | Painter         | 120000        |
| History          | Painter         | 50000         |
| Music            | Packard         | 80000         |
| Physics          | Watson          | 70000         |

CREATE TABLE course

```
(course_id    VARCHAR(7),  
 title        VARCHAR(50),  
 dept_name    VARCHAR(20),  
 credit       NUMERIC(2,0),  
 PRIMARY KEY (course_id),  
 FOREIGN KEY (dept_name) REFERENCES department );
```

# Editing Tables

DROP TABLE table\_name;

remove the table

DELETE FROM table\_name  
WHERE predicate;

delete tuples satisfying  
the predicate

ALTER TABLE table\_name  
ADD column type;

add a column

ALTER TABLE table\_name  
DROP column;

remove a column

# Inserting

In MySQL we can insert into a table with the command:

```
INSERT INTO table_name  
VALUES (value1, value2, ..., valuen);
```

OR

```
INSERT INTO table_name (col1, col2, ..., coln)  
VALUES (value1, value2, ..., valuen);
```

OR

```
INSERT INTO table_name  
SELECT QUERY
```

For example:

```
INSERT INTO instructor  
SELECT ID, name, dept_name, 18000  
FROM student  
WHERE dept_name = 'Music' AND tot_cred > 144;
```

# Updating

In MySQL we can update a table with the command:

```
UPDATE table_name  
    SET attribute = new_value
```

OR

```
UPDATE table_name  
    SET attribute = new_value  
    WHERE predicate or select statement;
```

OR

```
UPDATE table_name  
    SET attribute = CASE  
                        WHEN predicate THEN result  
                        WHEN predicate THEN result  
                        ...  
                        WHEN predicaten THEN resultn  
                        ELSE result0  
    END
```